

Towards a fully formalised proof of the syllepsis in triply monoidal ω -categories

Samuel Mimram, Eric Finster, Thibaut Benjamin

Ecole Polytechnique

07/07/2018

Goal

Give a formalisation of the definition of syllepsis in triply monoidal ω -categories, using the type theory CaTT defined by Finster and Mimram

The implementation we are using is available at <https://github.com/ThiBen/catt>

Motivations

- First proof for which the prover is needed

expected, but not proven result

Motivations

- First proof for which the prover is needed
- Extensive test of the implementation

various bugs could stay unnoticed on smaller projects

Motivations

- First proof for which the prover is needed
- Extensive test of the implementation
- Understanding of the practical use of CaTT

will guide the future developments

Outline

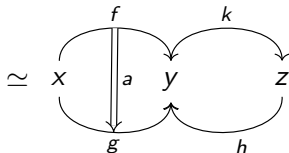
Type: globular sets

- To start with, we have a type theory with two type constructors

$$\frac{\Gamma \vdash}{\Gamma \vdash \star} \qquad \frac{\Gamma \vdash A \quad \Gamma \vdash x : A \quad \Gamma \vdash y : A}{\Gamma \vdash x \xrightarrow[A]{} y}$$

- Contexts of this theory describe globular sets in the following sense

$x : \star, y : \star, z : \star,$
 $f : x \rightarrow y, g : x \rightarrow y,$
 $h : z \rightarrow y, k : y \rightarrow z,$
 $a : f \rightarrow g$



Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors

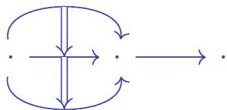
Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme

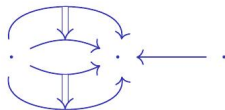
Definition

Pasting scheme are the globular sets that are in a "generic composable position"

Example



a pasting scheme



not a pasting scheme

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme

Examples

Operation	Pasting scheme	Result
Identity	x	$x \rightarrow x$

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme

Examples

Operation	Pasting scheme	Result
Identity	x	$x \rightarrow x$
Composition	$x \xrightarrow{f} y \xrightarrow{g} z$	$x \rightarrow z$

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme

Examples

Operation	Pasting scheme	Result
Identity	x	$x \rightarrow x$
Composition	$x \xrightarrow{f} y \xrightarrow{g} z$	$x \rightarrow z$
Associativity	$x \xrightarrow{f} y \xrightarrow{g} z \xrightarrow{h} w$	$f \star (g \star h) \rightarrow (f \star g) \star h$

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme

Examples

Operation	Pasting scheme	Result
• Identity	x	$x \rightarrow x$
Composition	$x \xrightarrow{f} y \xrightarrow{g} z$	$x \rightarrow z$
Associativity	$x \xrightarrow{f} y \xrightarrow{g} z \xrightarrow{h} w$	$f \star (g \star h) \rightarrow (f \star g) \star h$
coh id (x : *) : x -> x		

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme

Examples

Operation	Pasting scheme	Result
Identity	x	$x \rightarrow x$
• Composition	$x \xrightarrow{f} y \xrightarrow{g} z$	$x \rightarrow z$
Associativity	$x \xrightarrow{f} y \xrightarrow{g} z \xrightarrow{h} w$	$f \star (g \star h) \rightarrow (f \star g) \star h$
coh comp	$(x : *)$ $(y : *)$ $(f : x \rightarrow y)$ $(z : *)$ $(g : y \rightarrow z) :$ $x \rightarrow z$	

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme

Examples

Operation	Pasting scheme	Result
Identity	x	$x \rightarrow x$
Composition	$x \xrightarrow{f} y \xrightarrow{g} z$	$x \rightarrow z$
• Associativity	$x \xrightarrow{f} y \xrightarrow{g} z \xrightarrow{h} w$	$f \star (g \star h) \rightarrow (f \star g) \star h$

`coh assoc (x : *) (y : *) (f : x -> y)`
`(z : *) (g : y -> z)`
`(w : *) (h : z -> w) :`
`comp(x y f w (comp y z g w h)) ->`
`comp(x z (comp x y f z g) w h)`

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme
- Checking coherences

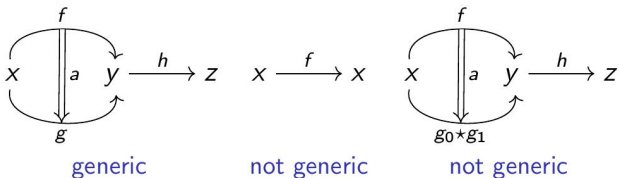
We add term constructor to ensure coherences correspond to correct operations of weak ω -category theory.

$$\frac{\Gamma \text{ pasting scheme} \quad \Gamma \vdash u \rightarrow v}{\Gamma \vdash \text{coh}_{\Gamma, u, v} : u \rightarrow v} \quad (+ \text{ extra conditions})$$

Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme
- Checking coherences
- Note: Genericity

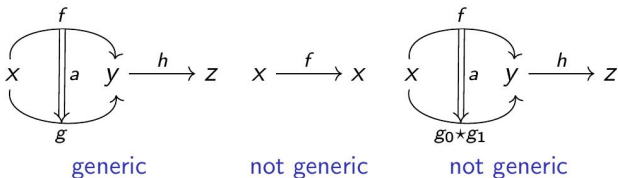
Example



Terms: weak ω -categories

- Operations and axioms of categories are generated using terms constructors
- Pasting scheme
- Coherence: operation of composing a pasting scheme
- Checking coherences
- Note: Genericity

Example



$\text{coh sq } (x : *) \text{ (f : } x \rightarrow x) : x \rightarrow x$

Some metatheoretical principles

- Implicit arguments

Example

Operation	Pasting scheme	Result
Associativity	$x \xrightarrow{f} y \xrightarrow{g} z \xrightarrow{h} w$	$f \star (g \star h) \rightarrow (f \star g) \star h$

```
coh assoc (x : *) (y : *) (f : x -> y)
           (z : *) (g : y -> z)
           (w : *) (h : z -> w) :
  comp(x y f w (comp y z g w h)) ->
  comp(x z (comp x y f z g) w h)
```

Some metatheoretical principles

- Implicit arguments

Example

Operation	Pasting scheme	Result
Associativity	$x \xrightarrow{f} y \xrightarrow{g} z \xrightarrow{h} w$	$f \star (g \star h) \rightarrow (f \star g) \star h$
coh assoc	$(x : *) (y : *) (z : *) (w : *)$ $(f : x \rightarrow y)$ $(g : y \rightarrow z)$ $(h : z \rightarrow w) :$	$\text{comp}(f (\text{comp } g \ h)) \rightarrow \text{comp}((\text{comp } f \ g) \ h)$

Some metatheoretical principles

- Implicit arguments
- Non-primitive operations

Example (squaring an endomorphism)

```
coh sq (x : *) (f : x -> x) : x -> x
```

Some metatheoretical principles

- Implicit arguments
- Non-primitive operations

Example (squaring an endomorphism)

```
coh sq (x : *) (f : x -> x) : x -> x
let sq (x : *) (f : x -> x) = comp f f
```

Some metatheoretical principles

- Implicit arguments
- Non-primitive operations
- Lifting of operations

Example

```
coh id0 (x : *) : x -> x
```

```
coh id1 (x : *) (y : *) (f : x -> y): f -> f
```

```
coh id2 (x: *) (y : *) (f : x -> y) (g : x ->  
y) (a : f -> g) : a -> a
```

Some metatheoretical principles

- Implicit arguments
- Non-primitive operations
- Lifting of operations

Solution

Declare only

```
coh id (x: *) : x -> x
```

and generate all the others on the fly

Outline

doubly monoidal ω -categories

Definition

A doubly monoidal ω -category is a weak ω -category with only one object, and one 1-cell

doubly monoidal ω -categories

Definition

A doubly monoidal ω -category is a weak ω -category with only one object, and one 1-cell

In CaTT:

- No way to specify that there is only one of a kind

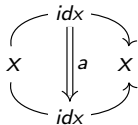
doubly monoidal ω -categories

Definition

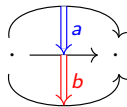
A doubly monoidal ω -category is a weak ω -category with only one object, and one 1-cell

In CaTT:

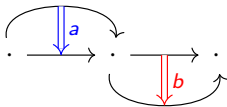
- No way to specify that there is only one of a kind
- Nevertheless, we can restrict ourselves to using only one object and one 1-cell



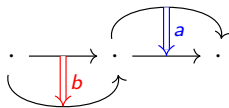
Definition of the braiding



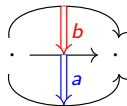
Definition of the braiding



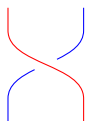
Definition of the braiding



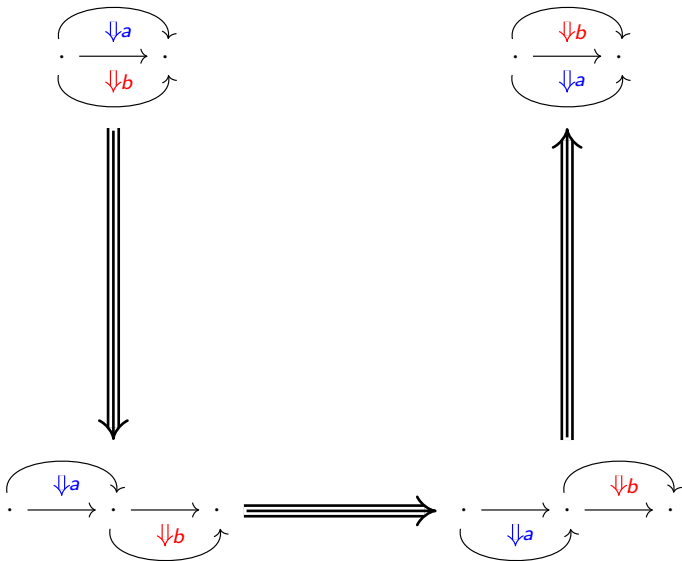
Definition of the braiding



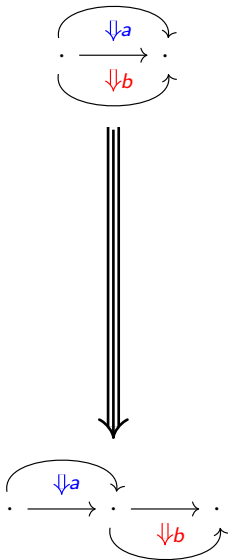
Definition of the braiding



Formalisation - First Plan

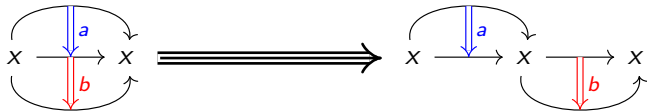


Formalisation - Attempt



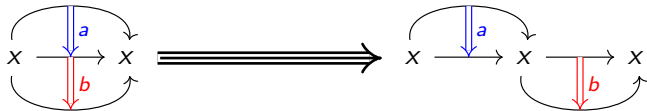
Formalisation Attempt

Translate these successive diagrams into terms of the theory



Formalisation Attempt

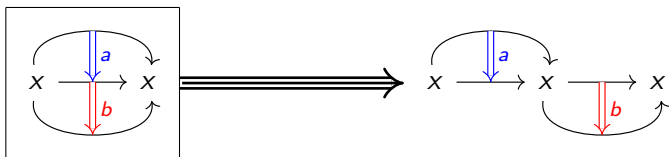
Translate these successive diagrams into terms of the theory



$\Gamma = (x : *) \quad (a : \text{id } x \rightarrow \text{id } x) \quad (b : \text{id } x \rightarrow \text{id } x)$

Formalisation Attempt

Translate these successive diagrams into terms of the theory

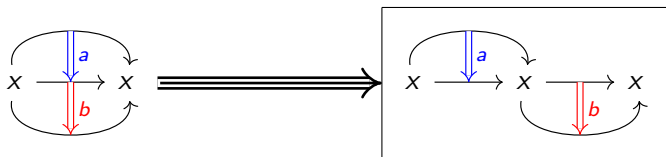


$\Gamma = (x : *) \quad (a : \text{id } x \rightarrow \text{id } x) \quad (b : \text{id } x \rightarrow \text{id } x)$

comp a b

Formalisation Attempt

Translate these successive diagrams into terms of the theory



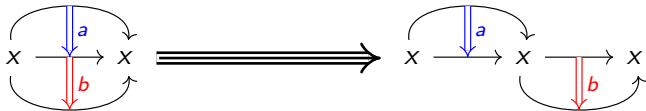
$\Gamma = (x : *) \quad (a : \text{id } x \rightarrow \text{id } x) \quad (b : \text{id } x \rightarrow \text{id } x)$

$\text{comp } a \ b$

$\text{comp } (\text{whiskR } a \ (\text{id } x)) \ (\text{whiskL } (\text{id } x) \ b)$

Formalisation Attempt

Translate these successive diagrams into terms of the theory



$\Gamma = (x : *) \ (a : \text{id } x \rightarrow \text{id } x) \ (b : \text{id } x \rightarrow \text{id } x)$

$\text{comp } a \ b \equiv \equiv \equiv \text{comp } (\text{whiskR } a \ (\text{id } x)) \ (\text{whiskL } (\text{id } x) \ b)$

Formalisation Attempt - First step

$\text{comp } a \ b \ \Longrightarrow \ \text{comp } (\text{whiskR } a \ (\text{id } x)) \ (\text{whiskL } (\text{id } x) \ b)$

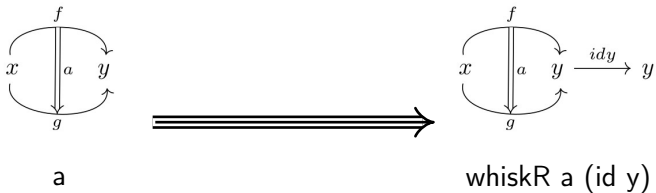
Formalisation Attempt - First step

a \Longrightarrow whiskR a (id x)

Idea

Handle separately a and b

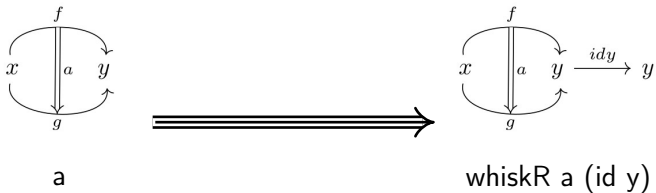
Formalisation Attempt - First step



To apply a coherence, we unfold into a generic position

$$\begin{aligned} (x : *) \quad (y : *) \quad (f : x \rightarrow y) \\ (g : x \rightarrow y) \quad (a : f \rightarrow g) \end{aligned}$$

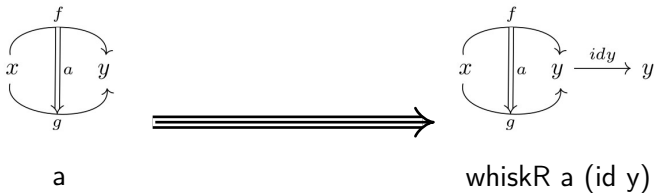
Formalisation Attempt - First step



Try a coherence

```
coh shiftL (x : *) (y : *) (f : x -> y)
           (g : x -> y) (a : f -> g) :
  a -> whiskR a (id y)
```

Formalisation Attempt - First step



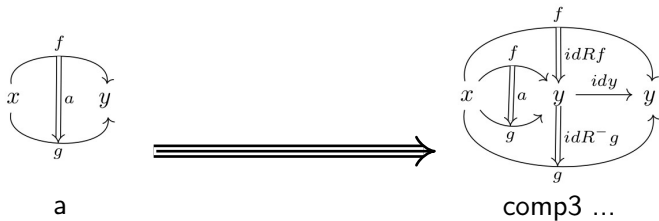
Result

This returns an error: the type of the result is ill-formed

$a : f \rightarrow g$

$\text{whiskR } a \text{ (id } y) : \text{comp } f \text{ (id } y) \rightarrow \text{comp } g \text{ (id } y)$

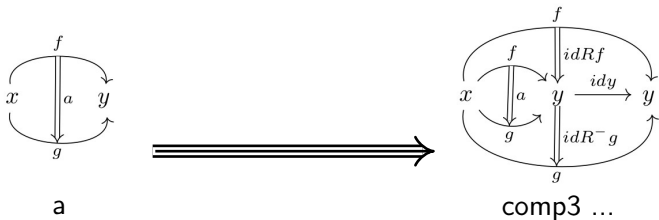
Formalisation Attempt - First step



Solution

Modify the target of the arrow

Formalisation Attempt - First step

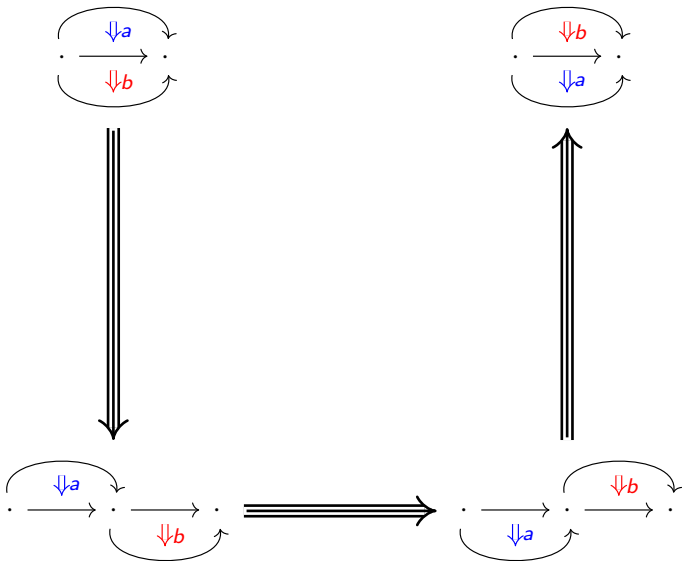


Solution

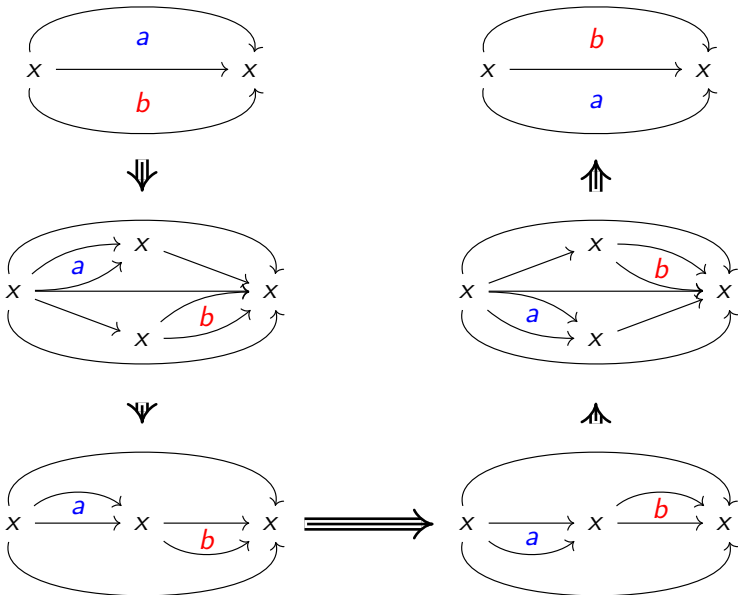
We now have a coherence

```
coh shiftL (x : *) (y : *) (f : x -> y)
           (g : x -> y) (a : f -> g) :
  a -> comp3 ...
```

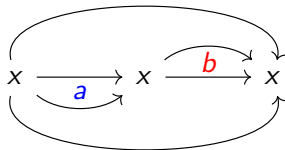
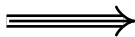
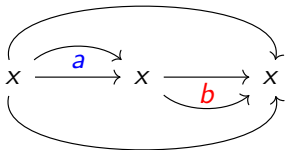
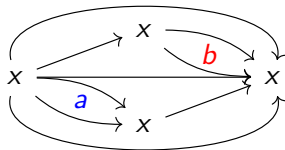
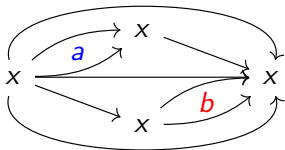
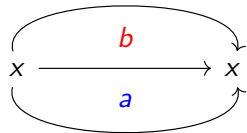
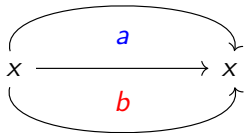
Formalisation - First Plan



Formalisation - Plan

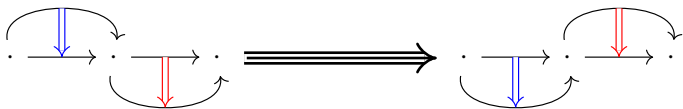


Formalisation - Core



Formalisation - Core

- Exchange law :



- This is a primitive operation of the theory

```
coh exch (x : *) (y : *) (f : x -> y) (g : x -> y)
      (a : f -> g)
      (z : *) (h : y -> z) (k : y -> z)
      (b : h -> k) :
  comp (whiskR a h) (whiskL g b) ->
  comp (whiskL f b) (whiskR a k)
```

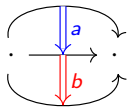
Outline

Inverse of the braiding

- The braiding has an obvious inverse

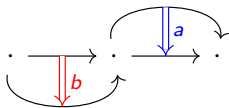
Inverse of the braiding

- The braiding has an obvious inverse
- Geometrically :



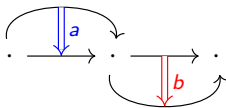
Inverse of the braiding

- The braiding has an obvious inverse
- Geometrically :



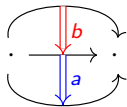
Inverse of the braiding

- The braiding has an obvious inverse
- Geometrically :



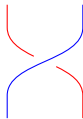
Inverse of the braiding

- The braiding has an obvious inverse
- Geometrically :



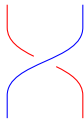
Inverse of the braiding

- The braiding has an obvious inverse
- Geometrically :



Inverse of the braiding

- The braiding has an obvious inverse
- Geometrically :



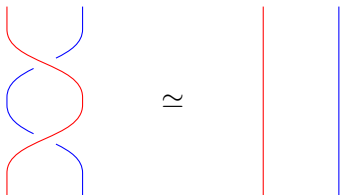
- Formally : The two definition have a very similar structure, but some of the arguments need to be inverted.

Cancellation witness

- The braiding and its inverse cancel out

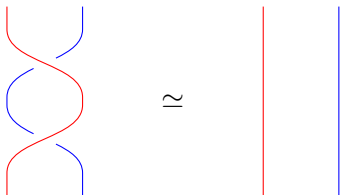
Cancellation witness

- The braiding and its inverse cancel out
- Geometrically :
Reidemeister move of type II



Cancellation witness

- The braiding and its inverse cancel out
- Geometrically :
Reidemeister move of type II



- Formally :
Providing cancellation witnesses for all intermediate, we can assemble them to build a cancellation for the complete braiding.

Automation

- This example shows how inversion works

Automation

- This example shows how inversion works
- It would be possible to generate inverses and cancellation witnesses algorithmically (but it is more technical than it looks, and has not been implemented yet)

Outline

triply monoidal ω -categories

Definition

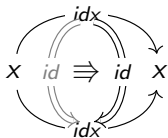
A triply monoidal ω -category is a weak ω -category with only one object, one 1-cell, and one 2-cell

triply monoidal ω -categories

Definition

A triply monoidal ω -category is a weak ω -category with only one object, one 1-cell, and one 2-cell

In CaTT:



triply monoidal ω -categories

Definition

A triply monoidal ω -category is a weak ω -category with only one object, one 1-cell, and one 2-cell

In CaTT:

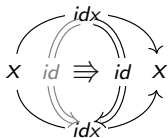
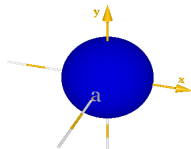
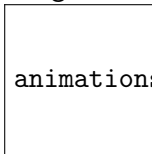


Diagram:



Geometric idea

- The "suspended" braiding



animations/eh1-3D/eh01.png